CS579 Computational Linguistics

Homework #4

Problem 1: Solve Exercise 4.2.1



Exercise 4.2.1 Give tableau proofs of the following formulas:

- 1. $\neg \neg p \rightarrow p$
- 2. $((p \rightarrow q) \rightarrow p) \rightarrow p$
- 3. $(\neg p \rightarrow \neg q) \rightarrow (q \rightarrow p)$
- 4. $p \to (p \land (q \lor p))$
- 5. $(p \lor (q \land r)) \rightarrow ((p \lor q) \land (p \lor r)).$

X See PDF page 115 of 1999 version of black book

Problem 2; Solve Exercise 4.4.5

Exercise 4.4.5 Give resolution proofs of the following formulas (recall that you were asked to convert the negations of these formulas into CNF in Exercise 4.4.1):

- 1. $\neg \neg p \rightarrow p$
- $2. ((p \to q) \to p) \to p$
 - 3. $(\neg p \rightarrow \neg q) \rightarrow (q \rightarrow p)$
- $4. \ p \to (p \land (q \lor p))$
- 5. $(p \lor (q \land r)) \rightarrow ((p \lor q) \land (p \lor r)).$

Compare these proofs with the tableau proofs requested in Exercise 4.2.1.

Exercise 5.6.2 Test the nonRedundantFactors/2 predicate. For example, test it on a list containing the single clause:

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[[a(m),a(Y),b(n,X),b(Y,Z),not(c(W)),not(c(f(Z)))]].
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Explain the result you get. Then try the next exercise.

Exercise 5.6.3 If you have tried the previous exercise, you will have noticed that nonRedundantFactors/2 over-generates new clauses. This is no surprise given its definition! Here is a suggestion to improve it by including a call to a predicate that discards clauses that are *subsumed* by others. (A clause C_1 subsumes a clause C_2 if and only if there is a substitution σ such that $C_1\sigma = C_2$. We say that C_2 is subsumed by C_1 .)

nonRedundantFactors([],[]).

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nonRedundantFactors([C1|L1],L5):-
findall(C2,nonRedFact(C1,C2),L3),
nonRedundantFactors(L1,L2),
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appendLists(L3,L2,L4), subsume(L4,L5).

Implement a subsumption check by giving a definition for subsume/2. Ensure you use unification with the occurs check.

See foResolution.pl in BB1 folder